

# High-Performance Scientific Computing

## Lecture 5: More OpenCL, MPI

MATH-GA 2011 / CSCI-GA 2945 · October 3, 2012

# Today

Tool of the day: Git

OpenCL: Device Language

OpenCL: Synchronization

Intro to MPI

# Bits and pieces

- HW1 grades sent
- HW2 graded soon
- HW3 due
- HW4 out tomorrow
- Cuda cluster accounts
- Mailing list messages

# Outline

Tool of the day: Git

OpenCL: Device Language

OpenCL: Synchronization

Intro to MPI

## Demo time

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## Demo time

# OpenCL Device Language

OpenCL device language is C99, with these differences:

- + Index getters
- + Memory space qualifiers
- + Vector data types
- + Many generic ('overloaded') math functions
- + Synchronization
- Recursion
- Fine-grained `malloc()`
- Function pointers





# Address Space Qualifiers

Type	Per	“Speed”
private*)	work item	super-fast
local	group	fast
global	grid	kinda slow

\*) default, so optional

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Should really discuss “speed” in terms of latency/bandwidth.

*Both* decrease with distance from the point of execution.

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**OpenCL: Synchronization**

Intro to MPI

# Concurrency and Synchronization

GPUs have layers of concurrency.

Each layer has its synchronization primitives.



# Concurrency and Synchronization

GPUs have layers of concurrency.

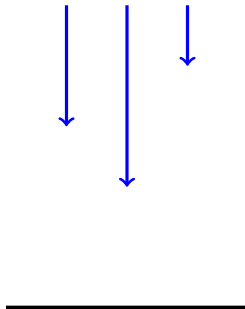
Each layer has its synchronization primitives.

- Intra-group:  
    `barrier(...)`,  
    `mem_fence(...)`  
    ... =  
    `CLK_{LOCAL,GLOBAL}_MEM_FENCE`
- Inter-group:  
    Kernel launch
- CPU-GPU:  
    Command queues



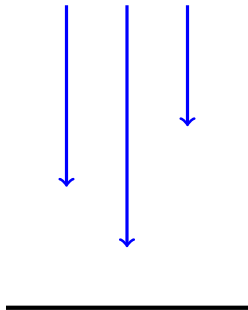
# Synchronization

What is a Barrier?



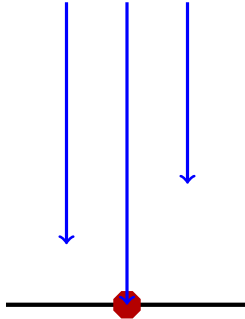
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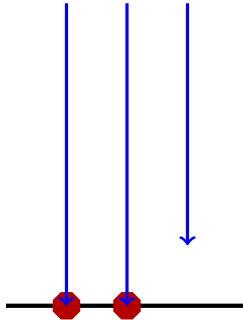
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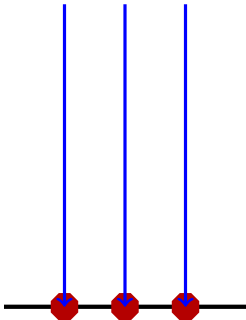
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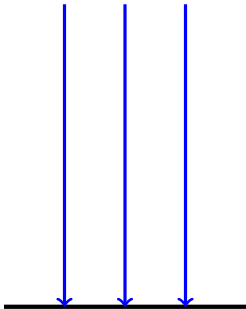
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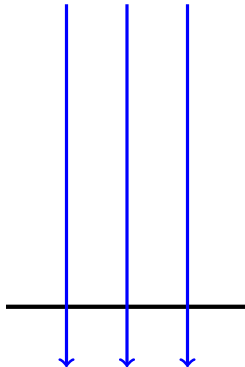
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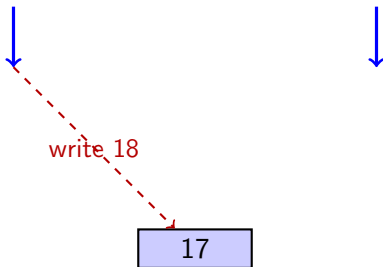
What is a Memory Fence?



17

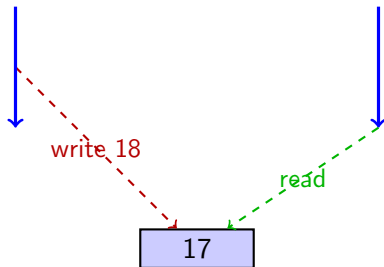
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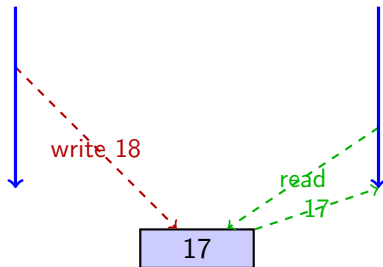
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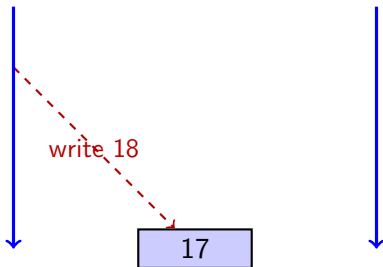
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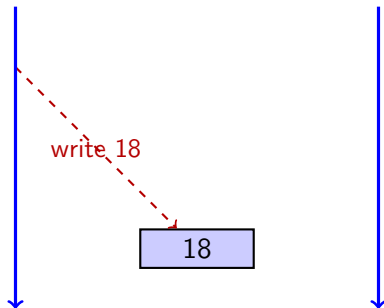
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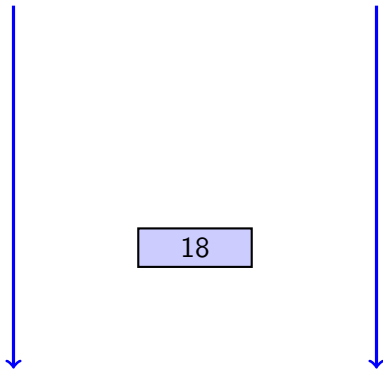
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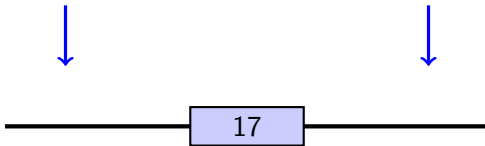
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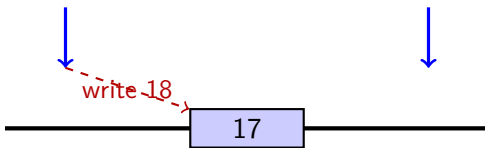
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What is a Memory Fence? An ordering restriction for memory access.



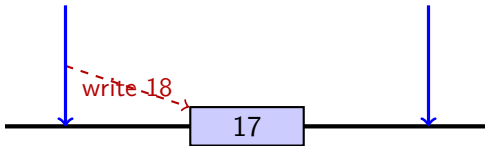
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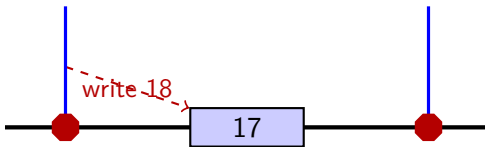
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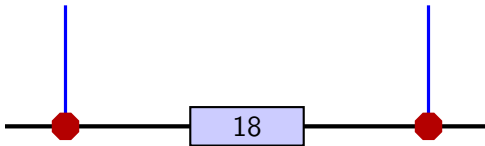
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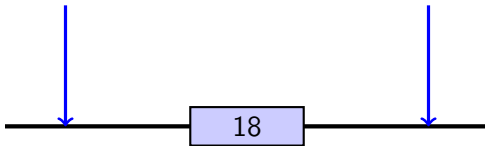
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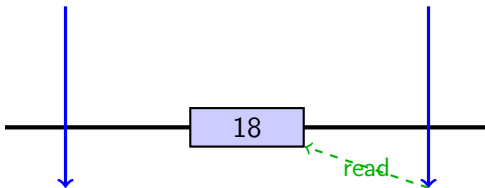
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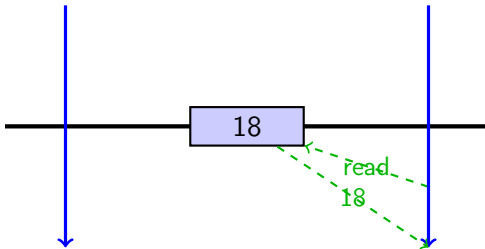
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# Synchronization between Groups

## Golden Rule:

Results of the algorithm must be independent of the order in which work groups are executed.

# Synchronization between Groups

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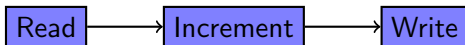
Results of the algorithm must be independent of the order in which work groups are executed.

## Consequences:

- Work groups may read the same information from global memory.
- But: Two work groups may not validly write different things to the same global memory.
- Kernel launch serves as
  - Global barrier
  - Global memory fence

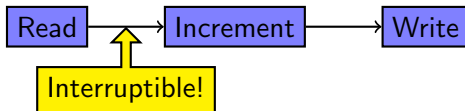
# Atomic Operations

Collaborative (inter-block) Global Memory Update:



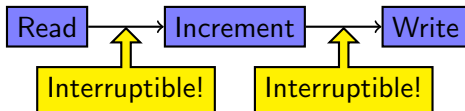
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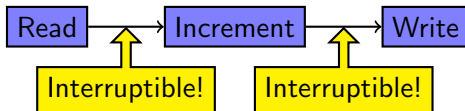
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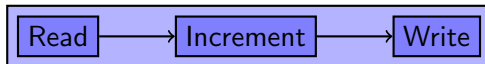


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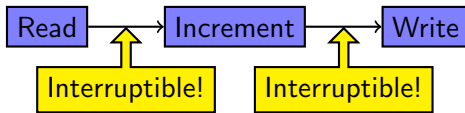


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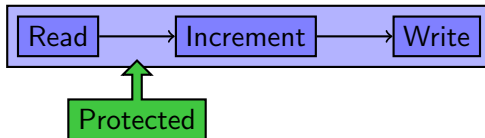


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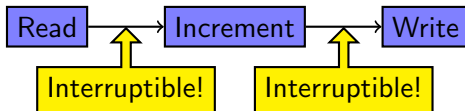


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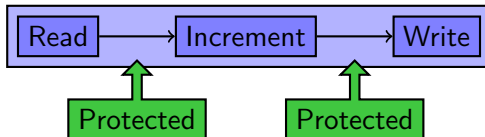


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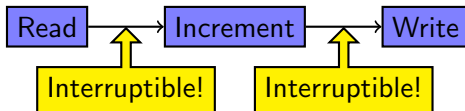


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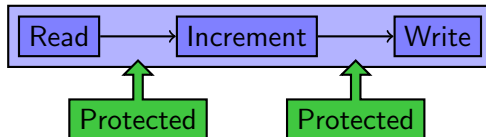


# Atomic Operations

Collaborative (inter-block) Global Memory Update:



Atomic Global Memory Update:



**How?**

```
atomic_{add,inc,cmpxchg,...}(int *global, int value);
```

# Atomic: Compare-and-swap

```
int atomic_cmpxchg (__global int *p, int cmp, int val)  
int atomic_cmpxchg (__local int *p, int cmp, int val)
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Does:

- Read the 32-bit value (referred to as old) stored at location pointed by p.
- Compute  $(old == cmp) ? val : old$ .
- Store result at location pointed by p.
- Returns old.

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Implement atomic float add?

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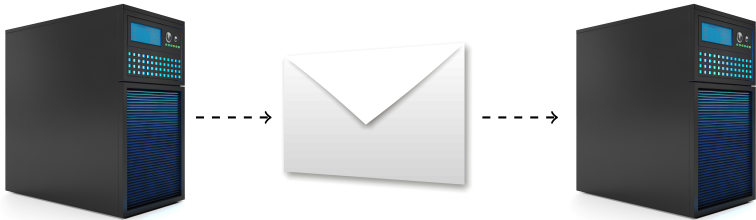
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OpenCL: Synchronization

**Intro to MPI**

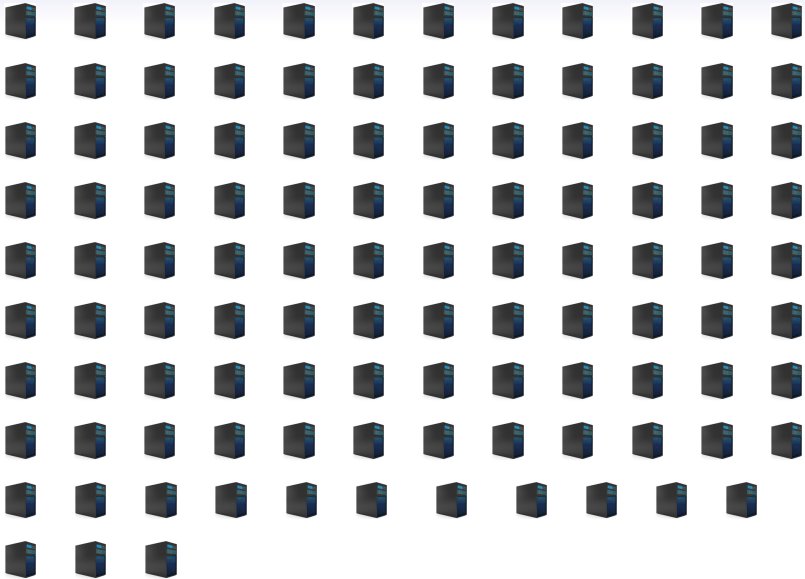
# MPI

## Message Passing Interface:

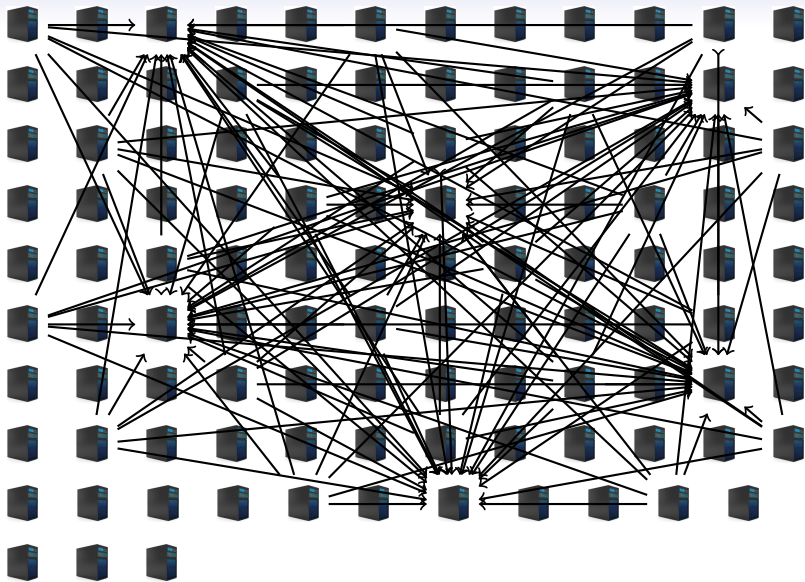




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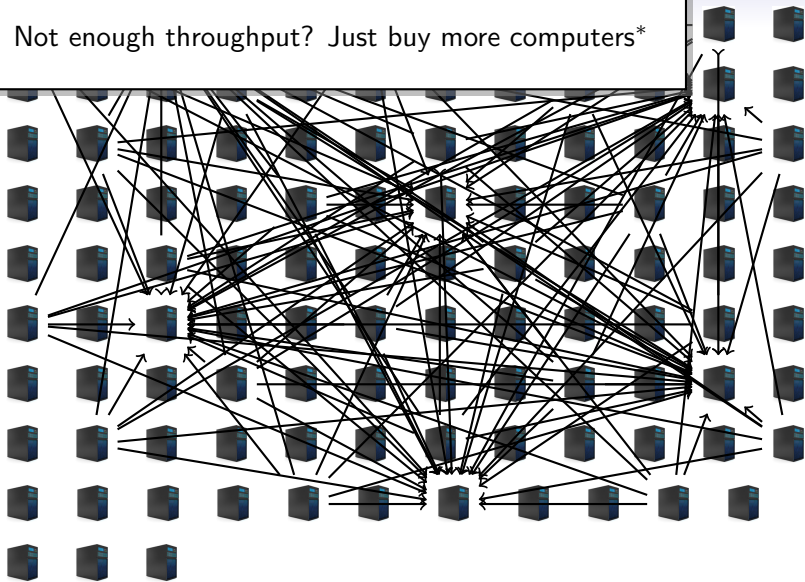


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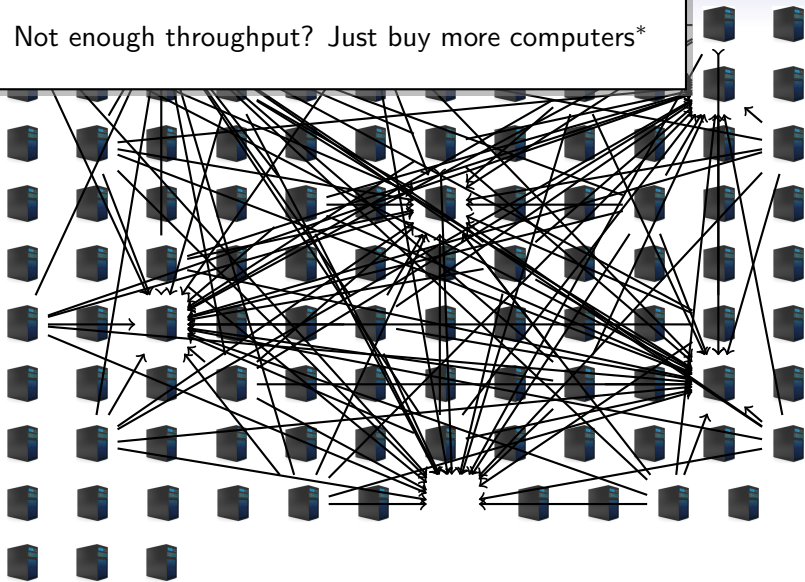
## MPI

Not enough throughput? Just buy more computers\*



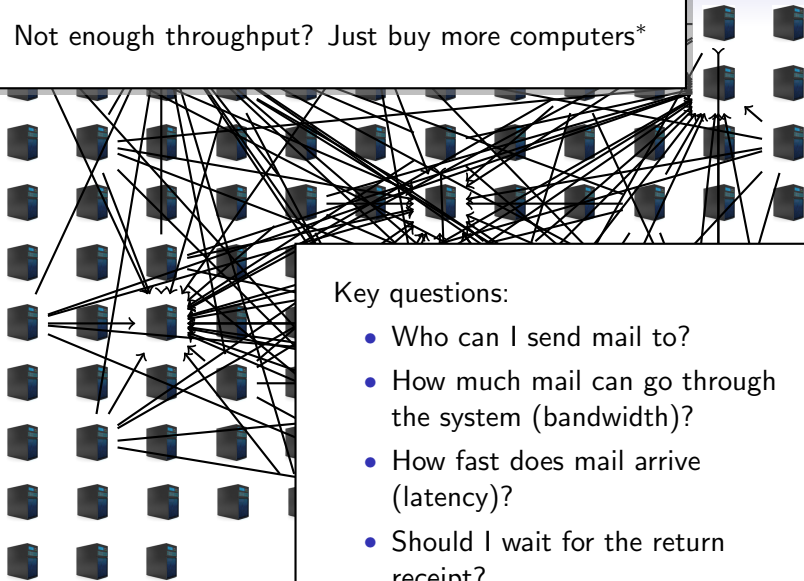
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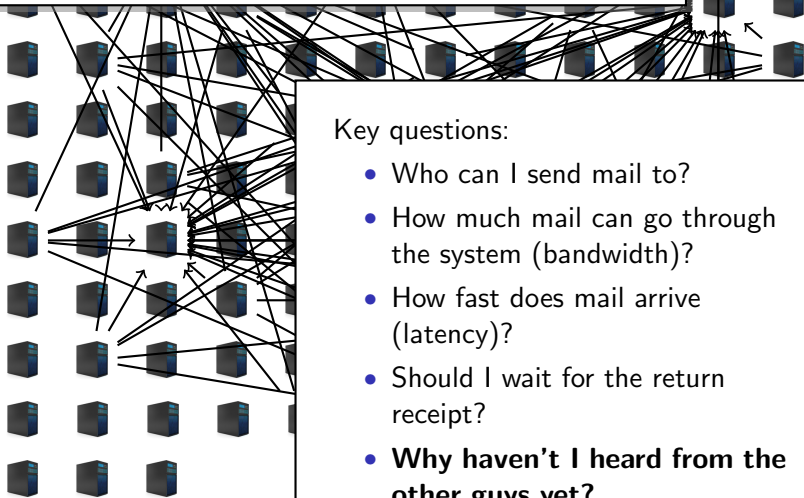
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Key questions:

- Who can I send mail to?
- How much mail can go through the system (bandwidth)?
- How fast does mail arrive (latency)?
- Should I wait for the return receipt?

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Key questions:

- Who can I send mail to?
- How much mail can go through the system (bandwidth)?
- How fast does mail arrive (latency)?
- Should I wait for the return receipt?
- **Why haven't I heard from the other guys yet?**

# MPI

## **MPI 3.0**

Born September 21, 2012

MPI 1.0: June 1994



## Demo time



# Questions?

?

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